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On Equivalency of Zero-divisor Codes via Classifying Their Idempotent Generator

Kai Lin Ong · Miin Huey Ang

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Abstract In 2009, Ted and Paul Hurley proposed a code construction method using group rings. These codes with single generator are termed *group ring codes* and in particular *zero-divisor codes* when using zero-divisors as generators. In this paper, we mainly study the equivalency of zero-divisor codes in F_2G having generator from $I(G)$, the set of all idempotents in F_2G . For abelian G , our previous notion of *generated idempotents* completely classified $I(G)$ by serving as its basis. Here, we first extend the notion of generated idempotents to study and classify some elements in $I(G)$ for non-abelian G . Later, the study is generally done on equivalency of zero-divisor codes in F_2G , then concentrating on those with idempotent generator. In particular, we affirm the conjecture “Every group ring code in F_2D_{2n} is equivalent to some in F_2C_{2n} ” in the cases where the generators are our classified idempotents. We also show that the equivalency of zero-divisor codes in F_2C_n with generated idempotent as generators can be established sufficiently on the generator property.

Keywords group ring code · zero-divisor code · idempotent · code equivalency

Mathematics Subject Classification (2010) 94B99 · 20C05 · 20D35

1 Introduction

The idea of defining codes as ideals in the group algebra F_qG for finite group G were introduced by Berman and MacWilliams independently [1,7]. These codes are termed *group codes*. In ring theory, a ring element whose square is itself is

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an *idempotent*. For decades, group codes having idempotent generators remain at the forefront of research. This is due to the fact that any group code C in F_qG with $\text{char}(F_q)|G$ can be expressed as $\bigoplus_{i=1}^l F_qGe_i$ for some $l \in \mathbb{Z}^+$ where each e_i is called a *primitive central idempotent generator* of C [8]. These primitive central idempotent generators played a crucial role in extracting the parameters and algebraic properties of the group codes [3,12].

In 2009, Hurley and Hurley [4] presented a brand new construction method of codes which involved group rings $RG = \{ \sum_{g \in G} a_g g | a_g \in R \}$.

Definition 1.1 Let RG be a group ring with $u \in RG$ and W be an R -submodule of RG . Define $f_{W,u} : W \rightarrow RG$ by $(w)f_{W,u} = wu$ for every $w \in W$. Then, $C = \text{Im}(f_{W,u})$ is called a *group ring code* in RG and u is called the *generator* of C .

The resultant group ring code has the property of being an R -submodule of RG . For every $w \in W$, w is known as a message word and wu is a codeword in $C = \text{Im}(f_{W,u})$. Clearly, $C = Wu = \{wu : w \in W\}$. For every codeword $c = \sum_{g \in G} a_g g \in C$, define its *support* $\text{supp}(c) = \{g \in G | a_g \neq 0\}$ and weight $wt(c) = |\text{supp}(c)|$. Also, it can be easily verified that $f_{W,u}$ is an R -linear map and thus C is an R -submodule of RG . Research work on group ring codes further led to the dichotomy of zero-divisor and unit-derived codes in RG , depending on the nature of their generator. From the point of zero-divisor codes, note that when $W = RG$, the resultant zero-divisor codes in RG possess ideal structure, thus coinciding with the notion of group codes. Such codes in the form of RGu are called ideal zero-divisor codes in RG . Various works on zero-divisor and unit-derived codes were established in the recent ten years, predominantly on the self-duality of codes and the construction of new convolutional codes [5,6].

The main study of equivalent codes is to identify the criteria that helps to classify codes that look different as a set but perform alike as codes from the theoretical and practical points. Theoretically, two codes are considered different if they exhibit different algebraic structures. Practically, two codes are considered different if they have different parameter sets.

In 1962, MacWilliams proved that two linear codes of the same length are equivalent if and only if there is a weight-preserving isomorphism between them [2]. A definition of equivalent zero-divisor codes in F_2G is given by modifying MacWilliams' result into the following group ring version.

Definition 1.2 Let G_1 and G_2 be two finite groups with $|G_1| = |G_2|$. Let W_1u_1 and W_2u_2 be zero-divisor codes in F_2G_1 and F_2G_2 respectively. Then, W_1u_1 and W_2u_2 are said to be equivalent if there exists an F_2 -linear isomorphism $\varphi : W_1u_1 \rightarrow W_2u_2$ such that for every codeword $c \in W_1u_1$, $wt(c) = wt((c)\varphi)$.

In general, despite the exhibition of different algebraic properties of the non-isomorphic G_1 and G_2 , there are 9 examples for which equivalency holds between zero-divisor codes in F_2D_{2n} and those in F_2C_{2n} , where D_{2n} and C_{2n} are the dihedral and cyclic groups of order $2n$ respectively [13]. This led to the conjecture "Every group ring code in F_2D_{2n} is equivalent to some group ring codes in F_2C_{2n} " [13]. As any arbitrary idempotent in F_2G can be chosen as a generator of a zero-divisor code $C = Wu$ in F_2G , the primary step to study equivalency is to

establish a way to classify and identify idempotents in F_2G . Let $I(G)$ be the set of all idempotents in F_2G . For abelian G , a new classification of idempotents that completely classified and identified all elements in $I(G)$ were introduced by Ong and Ang in 2017 [9,10]. This was done by first affirming the existence of a proper subset of $I(G)$ where the support of each of its element is generated by an element in G .

Definition 1.3 *Let G be an abelian group and let $e \in I(G)$ with $|\text{supp}(e)| = k$. If $e = \sum_{i=0}^{k-1} g^{2^i}$ for $g \in \text{supp}(e)$, then e is said to be a generated idempotent in F_2G with generator g and is denoted by $\langle g \rangle_{Id}$.*

Throughout this paper, $I_{\langle Id \rangle}(G)$ denotes the set of all generated idempotents in F_2G . We have shown that apart from a subset of $I(G)$, $I_{\langle Id \rangle}(G)$ is precisely a basis for $I(G)$. The readers can refer to Corollary 3.6 in [9].

Theorem 1.4 *Let G be an abelian group. Then, $I_{\langle Id \rangle}(G)$ is a basis for $I(G)$.*

The identification of $I_{\langle Id \rangle}(G)$ for abelian G began with the following theorem, which in fact holds generally for any G .

Theorem 1.5 *Let $g \in G$. Then, $\langle g \rangle_{Id}$ is a non-trivial generated idempotent in F_2G if and only if $\text{ord}(g)$ is odd with $\text{ord}(g) > 1$.*

Readers who are interested with the details of the complete identification of $I_{\langle Id \rangle}(G)$ for abelian G are referred to [9,10]. We include a concise example to illustrate the identification of $I(G)$ for an abelian G . Note that, throughout this paper, C_n denotes a cyclic group of order n whereas $C_{\langle g \rangle}$ denotes a cyclic group with generator g .

Example 1.6 *Consider F_2C_{14} with $C_{14} = C_{\langle x \rangle}$. The set of all generated idempotents in F_2C_{14} comprises three elements, that is $I_{\langle Id \rangle}(G) = \{1, \langle x^2 \rangle_{Id} = x^2 + x^4 + x^8, \langle x^6 \rangle_{Id} = x^6 + x^{10} + x^{12}\}$. Hence, by Theorem 1.4, we have $I(C_{14}) = L_{F_2}(I_{\langle Id \rangle}(C_{14}))$, the linear span of $I_{\langle Id \rangle}(C_{14})$ over F_2 with $|I(C_{14})| = 2^3 = 8$.*

From [9], it is instructive to note that we had shown for distinct $\langle g \rangle_{Id}, \langle h \rangle_{Id} \in I_{\langle Id \rangle}(G)$, $\text{supp}\langle g \rangle_{Id} \cap \text{supp}\langle h \rangle_{Id} = \emptyset$. This eventually led to the following classification of elements in $I(G)$ for abelian G .

Theorem 1.7 *Every non-zero $e \in I(G)$ is either a generated idempotent or a finite sum of some generated idempotents in F_2G .*

Unfortunately, for non-abelian G , the extension of this result fails to hold as shown by a counterexample in the dihedral group ring over $D_6 = \langle a, b | a^3 = b^2 = 1, aba = b^{-1} \rangle$; the element $a + ab + a^2b \in F_2D_6$ belongs to neither of the forms. However, we persist to focus on those idempotents in the form of Theorem 1.7 as Theorem 1.5 convincingly suggests that there is an abundance of $e \in I_{\langle Id \rangle}(G)$ for most of the groups G .

The paper is organized as follow. In Section 2, as for non-abelian G , $I(G)$ no longer possesses vector space structure, and so the study on $I(G)$ is then primarily focused on identifying and studying the properties of a subset of $I_{\langle Id \rangle}(G)$, called the set of joint idempotents, that is, $J(G) = \{\langle g \rangle_{Id} \in I_{\langle Id \rangle}(G) | \langle g \rangle_{Id} + \langle h \rangle_{Id} \in I(G), \text{ for every } \langle h \rangle_{Id} \in I_{\langle Id \rangle}(G)\}$. A complete classification of $J(G)$ for several

families of non-abelian F_2G is given in Section 3. Section 4 consists of a comprehensive study on the conditions that affirms the equivalency between zero-divisor codes in F_2G . The results are then utilized in Section 5, in which we will show that every zero-divisor code in F_2D_{2n} with idempotent generator in either form mentioned in Theorem 1.7 has an equivalent form in F_2C_{2n} . Also, we establish the equivalency of zero-divisor codes in F_2C_n with generated idempotent as generator by solely looking at the generator property.

2 Idempotents in Non-abelian F_2G

Throughout this section, G is fixed to be a non-abelian group. Consider distinct generated idempotents $\langle g \rangle_{Id}, \langle h \rangle_{Id} \in I_{\langle Id \rangle}(G)$. In general, Theorem 1.4 for general G does not hold as $\langle g \rangle_{Id} + \langle h \rangle_{Id}$ is not necessarily in $I(G)$. This can be seen from the counterexample $G = A_4$, the alternating group on $\{1, 2, 3, 4\}$. It can be verified that $\langle (123) \rangle_{Id} = (123) + (132)$, $\langle (124) \rangle_{Id} = (124) + (142) \in I_{\langle Id \rangle}(G)$ but $\langle (123) \rangle_{Id} + \langle (124) \rangle_{Id} \notin I(G)$. This triggers the introduction of the notion of joinability of generated idempotents.

Definition 2.1 For distinct $\langle g \rangle_{Id}, \langle h \rangle_{Id} \in I_{\langle Id \rangle}(G)$, $\langle g \rangle_{Id}$ is said to be joinable with $\langle h \rangle_{Id}$ if $\langle g \rangle_{Id} + \langle h \rangle_{Id} \in I(G)$.

For distinct $\langle g \rangle_{Id}, \langle h \rangle_{Id} \in I_{\langle Id \rangle}(G)$, note that $(\langle g \rangle_{Id} + \langle h \rangle_{Id})^2 = \langle g \rangle_{Id} + \langle h \rangle_{Id}$ if and only if $\langle g \rangle_{Id}\langle h \rangle_{Id} + \langle h \rangle_{Id}\langle g \rangle_{Id} = 0$. This results in the following bi-condition, which will be frequently utilized subsequently in this section: $\langle g \rangle_{Id}$ is joinable with $\langle h \rangle_{Id}$ if and only if $\langle g \rangle_{Id}\langle h \rangle_{Id} = \langle h \rangle_{Id}\langle g \rangle_{Id}$. The next result ensures that if $\langle g \rangle_{Id}$ is joinable with $\langle h \rangle_{Id}$, then $C_{\langle g \rangle}$ must permute with $C_{\langle h \rangle}$.

Proposition 2.2 For distinct $\langle g \rangle_{Id}, \langle h \rangle_{Id} \in I_{\langle Id \rangle}(G)$, if $\langle g \rangle_{Id}$ is joinable with $\langle h \rangle_{Id}$, then $C_{\langle g \rangle}C_{\langle h \rangle} = C_{\langle h \rangle}C_{\langle g \rangle}$.

Proof Suppose that $\langle g \rangle_{Id}$ is joinable with $\langle h \rangle_{Id}$, then $\langle g \rangle_{Id}\langle h \rangle_{Id} = \langle h \rangle_{Id}\langle g \rangle_{Id}$. Then, for every $i \in \{0, 1, 2, \dots, |\text{supp}\langle g \rangle_{Id}| - 1\}$ and $j \in \{0, 1, 2, \dots, |\text{supp}\langle h \rangle_{Id}| - 1\}$, there exist $l \in \{0, 1, 2, \dots, |\text{supp}\langle h \rangle_{Id}| - 1\}$ and $m \in \{0, 1, 2, \dots, |\text{supp}\langle g \rangle_{Id}| - 1\}$ such that

$$g^{2^i}h^{2^j} = h^{2^l}g^{2^m}. \quad (2.1)$$

Note that for each $s \in \mathbb{Z}_n^+$, there exists unique $\{x_{i_1}, x_{i_2}, \dots, x_{i_r}\} \subset \mathbb{N}$, such that $s = 2^{x_{i_1}} + 2^{x_{i_2}} + \dots + 2^{x_{i_r}}$. Let $g^\alpha h^\beta \in C_{\langle g \rangle}C_{\langle h \rangle}$, this implies that there exists unique $\{i_1, i_2, \dots, i_r\}, \{j_1, j_2, \dots, j_s\} \subset \mathbb{N}$ such that $g^\alpha h^\beta = g^{\sum_{k=1}^r 2^{i_k}} h^{\sum_{k=1}^s 2^{j_k}} = \prod_{k=1}^r g^{2^{i_k}} \prod_{k=1}^s h^{2^{j_k}}$. Then, by (2.1), note that $\prod_{k=1}^r g^{2^{i_k}} \prod_{k=1}^s h^{2^{j_k}} = \prod_{k=1}^s h^{2^{l_k}} \prod_{k=1}^r g^{2^{m_k}} = h^{\sum_{k=1}^s 2^{l_k}} g^{\sum_{k=1}^r 2^{m_k}}$. Thus, $g^\alpha h^\beta = h^{\sum_{k=1}^s 2^{l_k}} g^{\sum_{k=1}^r 2^{m_k}} \in C_{\langle h \rangle}C_{\langle g \rangle}$ and this results in $C_{\langle g \rangle}C_{\langle h \rangle} \subseteq C_{\langle h \rangle}C_{\langle g \rangle}$. By a similar argument, it can be shown that $C_{\langle h \rangle}C_{\langle g \rangle} \subseteq C_{\langle g \rangle}C_{\langle h \rangle}$ and thus $C_{\langle g \rangle}C_{\langle h \rangle} = C_{\langle h \rangle}C_{\langle g \rangle}$. \square

We term those $\langle g \rangle_{Id} \in I_{\langle Id \rangle}(G)$ which are joinable with every generated idempotent in $I_{\langle Id \rangle}(G)$ as follows.

Definition 2.3 Let $\langle g \rangle_{Id} \in I_{\langle Id \rangle}(G)$. Then, $\langle g \rangle_{Id}$ is termed a joint idempotent if for every $\langle h \rangle_{Id} \in I_{\langle Id \rangle}(G)$, $\langle g \rangle_{Id}$ is joinable with $\langle h \rangle_{Id}$.

Throughout this paper, the set of all joint idempotents in F_2G is denoted as $J(G)$. Then, $\langle g \rangle_{Id} \in J(G)$ if and only if for every $\langle h \rangle_{Id} \in I_{(Id)}(G)$, $\langle g \rangle_{Id} \langle h \rangle_{Id} = \langle h \rangle_{Id} \langle g \rangle_{Id}$. For every G , $J(G)$ is always non-empty as the trivial generated idempotent $1 \in J(G)$. Moreover, we prove that $\{\langle g \rangle_{Id} \in I_{(Id)}(G) | g \in Z(G)\} \subseteq J(G)$, where $Z(G)$ denotes the centre of G .

Proposition 2.4 *Let $\langle g \rangle_{Id} \in I_{(Id)}(G)$ with $g \in Z(G)$. Then, $\langle g \rangle_{Id} \in J(G)$.*

Proof Let $\langle h \rangle_{Id} \in I_{(Id)}(G)$. Since $Z(G)$ is a subgroup of G , then $\text{supp}\langle g \rangle_{Id} \subseteq Z(G)$. This results in $\langle g \rangle_{Id} \langle h \rangle_{Id} = \langle h \rangle_{Id} \langle g \rangle_{Id}$. \square

Recall that Proposition 2.2 reflects that joinability of $\langle g \rangle_{Id}, \langle h \rangle_{Id} \in I_{(Id)}(G)$ results in their corresponding cyclic subgroups $C_{\langle g \rangle}$ and $C_{\langle h \rangle}$ of G to be permuted. The next proposition gives a necessary condition for $\langle g \rangle_{Id} \in J(G)$.

Proposition 2.5 *Let $\langle g \rangle_{Id} \in I_{(Id)}(G)$. If $\langle g \rangle_{Id} \in J(G)$, then $C_{\langle g \rangle}$ permutes with every odd order subgroup of G .*

Proof It follows from Theorem 1.5 and Proposition 2.2 that $C_{\langle g \rangle}$ permutes with every odd order cyclic subgroup of G , that is $C_{\langle g \rangle} C_{\langle h \rangle} = C_{\langle h \rangle} C_{\langle g \rangle}$ for every odd order cyclic subgroup $C_{\langle h \rangle}$ of G . Let K be an odd order subgroup of G . For every $g^i k \in C_{\langle g \rangle} K$, $C_{\langle g \rangle} C_{\langle k \rangle} = C_{\langle k \rangle} C_{\langle g \rangle}$ results in $g^i k = k^l g^j \in C_{\langle k \rangle} C_{\langle g \rangle} \subseteq KC_{\langle g \rangle}$ for some $k^l \in C_{\langle k \rangle}$ and $g^j \in C_{\langle g \rangle}$. Thus, $C_{\langle g \rangle} K \subseteq KC_{\langle g \rangle}$. Using a similar argument, it can be verified that $KC_{\langle g \rangle} \subseteq C_{\langle g \rangle} K$, hence $C_{\langle g \rangle} K = KC_{\langle g \rangle}$. \square

Proposition 2.5 indicates the need of the notion of odd-permutable subgroups as defined next.

Definition 2.6 *Let H be an odd order subgroup of G . Then, H is termed an odd-permutable subgroup of G if H permutes with every odd order subgroup of G .*

In group theory, the permutability of a subgroup always guarantee its subnormality [11]. In conjunction with the odd-permutability of a subgroup, we introduce the notion of odd-subnormality.

Definition 2.7 *Let H be an odd order subgroup of G . Then, H is said to be an odd-subnormal subgroup of G if there exists $1 = G_0 \triangleleft G_1 \triangleleft G_2 \triangleleft \cdots \triangleleft G_r$, a subnormal series where $|G_i|$ is odd for every $i \in \{1, 2, \dots, r-1\}$ with $H = G_j$ for some $j \in \{0, 1, 2, \dots, r-1\}$.*

Let H and K be distinct odd order subgroups of G . It follows easily that HK and any conjugate of H must have odd order. Then, with these assertions, the next theorem directly follows from the context in Section 13.2 in [11].

Theorem 2.8 *Every odd-permutable subgroup of G is an odd-subnormal subgroup of G .*

We thus obtain another necessary condition for $\langle g \rangle_{Id} \in J(G)$ as follows.

Theorem 2.9 *Let $\langle g \rangle_{Id} \in I_{(Id)}(G)$. If $\langle g \rangle_{Id} \in J(G)$, then $C_{\langle g \rangle}$ is odd-subnormal in G .*

Proof This follows from Proposition 2.5 and Theorem 2.8. \square

3 Joint Idempotents in Some Non-abelian F_2G

In this section, $J(G)$ of three families of F_2G are studied. The first two families involve groups which are semidirect product of their two cyclic subgroups.

Definition 3.1 Let C_m and C_n be two non-trivial cyclic subgroups of G . Then, G is said to be a semidirect product of C_m by C_n , denote as $G = C_m \rtimes C_n$, if the following conditions are satisfied:

1. $G = C_m C_n = \{xy | x \in C_m, y \in C_n\}$ with C_m is normal in G ,
2. $C_m \cap C_n = \{1\}$.

3.1 Non-abelian Groups of Order pq

Let G be of order pq , where p and q are two distinct odd primes, with $p > q$. Note that $G = C_p \rtimes C_q$. Therefore, $G = \langle x, y | x^p = y^q = 1, xy = yx^\alpha \rangle$ for some $\alpha \in \mathbb{Z}_p^* \setminus \{1\}$. Each element in G has a general form of $x^i y^j$, where $i \in \{0, 1, 2, \dots, p-1\}$ and $j \in \{0, 1, 2, \dots, q-1\}$. By Theorem 1.5, we have $I_{\langle Id \rangle}(G) = \{\langle x^i y^j \rangle_{Id} | i \in \{0, 1, 2, \dots, p-1\}, j \in \{0, 1, 2, \dots, q-1\}\}$. Then, $I_{\langle Id \rangle}(G)$ can be partitioned into $\{\langle x^i y^j \rangle_{Id} | j \neq 0\}$ and $\{\langle x^i y^j \rangle_{Id} | j = 0\}$. We first show that for the former set, $\{\langle x^i y^j \rangle_{Id} | j \neq 0\} \cap J(G) = \emptyset$.

Proposition 3.2 Let p and q be two distinct odd primes, with $p > q$ and $G = C_p \rtimes C_q = \langle x, y | x^p = y^q = 1, xy = yx^\alpha \rangle$ for some $\alpha \in \mathbb{Z}_p^* \setminus \{1\}$. Then, for any $\langle x^{k_1} y^{k_2} \rangle_{Id} \neq \langle x^i y^j \rangle_{Id}$ with $j \neq 0$ and $k_2 \neq 0$, $\langle x^{k_1} y^{k_2} \rangle_{Id}$ is not joinable with $\langle x^i y^j \rangle_{Id}$ if and only if $C_{\langle x^{k_1} y^{k_2} \rangle} \neq C_{\langle x^i y^j \rangle}$.

Proof Note that every proper subgroup of G is of order either p or q . Clearly, C_p is a Sylow p -subgroup of G . As C_p is normal in G , C_p is the unique Sylow p -subgroup of G . Thus, $C_{\langle x^i y^j \rangle}$ must be a cyclic subgroup of G with order q for every $j \neq 0$.

Suppose that $C_{\langle x^{k_1} y^{k_2} \rangle} = C_{\langle x^i y^j \rangle}$, then clearly $\langle x^{k_1} y^{k_2} \rangle_{Id}$ is joinable with $\langle x^i y^j \rangle_{Id}$ as they are both in $I_{\langle Id \rangle}(C_{\langle x^i y^j \rangle})$. Conversely, suppose that $C_{\langle x^{k_1} y^{k_2} \rangle} \neq C_{\langle x^i y^j \rangle}$. Then, both $C_{\langle x^{k_1} y^{k_2} \rangle}$ and $C_{\langle x^i y^j \rangle}$ are of order q with $C_{\langle x^{k_1} y^{k_2} \rangle} \cap C_{\langle x^i y^j \rangle} = \{1\}$. If $C_{\langle x^{k_1} y^{k_2} \rangle}$ permutes with $C_{\langle x^i y^j \rangle}$, then it follows that $C_{\langle x^{k_1} y^{k_2} \rangle} C_{\langle x^i y^j \rangle}$ is a subgroup of G with $|C_{\langle x^{k_1} y^{k_2} \rangle} C_{\langle x^i y^j \rangle}| = q^2$. This contradicts Lagrange's Theorem. Thus, $C_{\langle x^{k_1} y^{k_2} \rangle}$ does not permute with $C_{\langle x^i y^j \rangle}$, Proposition 2.2 then concludes that $\langle x^{k_1} y^{k_2} \rangle_{Id}$ is not joinable with $\langle x^i y^j \rangle_{Id}$. \square

The following corollary is a direct result from Proposition 3.2.

Corollary 3.3 Let p and q be two distinct odd primes, with $p > q$ and $G = C_p \rtimes C_q = \langle x, y | x^p = y^q = 1, xy = yx^\alpha \rangle$ for some $\alpha \in \mathbb{Z}_p^* \setminus \{1\}$. Then, $\langle x^i y^j \rangle_{Id} \notin J(G)$ if $j \neq 0$.

Let us introduce a useful notation here. For $s \in \{1, 2, \dots, p-1\}$, the 2-cyclotomic coset of \mathbb{Z}_p containing s is defined as $\mathfrak{C}_s = \{s(2^k) | k \in \mathbb{N}\}$. For the remaining $\{\langle x^i y^j \rangle_{Id} | j = 0\}$, the value of $\alpha \in \mathbb{Z}_p^* \setminus \{1\}$ from $G = C_p \rtimes C_q = \langle x, y | x^p = y^q = 1, xy = yx^\alpha \rangle$ determines their joinability. Before that, the following technical lemma is needed.

Lemma 3.4 Let $m, n \in \mathbb{Z}^+$ and $s \in \{1, 2, \dots, m-1\}$. Let $G = \langle x, y | x^m = y^n = 1, xy = yx^\alpha \rangle$ for some $\alpha \in \mathbb{Z}_m^* \setminus \{1\}$. For every $i \in \{0, 1, 2, \dots, m-1\}$ and $j \in \{0, 1, 2, \dots, n-1\}$, $x^{s2^{k_1}}(x^i y^j)^{2^{k_2}} = (x^i y^j)^{2^{k_2}} x^{s2^{k_1} \alpha^{2^{k_2} j}}$ for $k_1, k_2 \in \mathbb{N}$.

Proof The proof can be directly shown using multidimensional mathematical induction on $k_1, k_2 \in \mathbb{N}$. \square

Theorem 3.5 Let p and q be two distinct odd primes, with $p > q$ and $G = C_p \rtimes C_q = \langle x, y | x^p = y^q = 1, xy = yx^\alpha \rangle$ for some $\alpha \in \mathbb{Z}_p^* \setminus \{1\}$. Then, for $s \in \{1, 2, \dots, p-1\}$, $\langle x^s \rangle_{Id} \in J(G)$ if $\alpha \in \mathfrak{C}_1$, where \mathfrak{C}_1 is the 2-cyclotomic coset of \mathbb{Z}_p containing 1.

Proof Note that $\langle x^s \rangle_{Id} \in J(G)$ if and only if $\langle x^s \rangle_{Id} \langle x^i y^j \rangle_{Id} = \langle x^i y^j \rangle_{Id} \langle x^s \rangle_{Id}$ for each pair of $i \in \{0, 1, 2, \dots, p-1\}$ and $j \in \{0, 1, 2, \dots, q-1\}$. Let $g = x^{s2^{k_1}}(x^i y^j)^{2^{k_2}} \in \text{supp}(\langle x^s \rangle_{Id} \langle x^i y^j \rangle_{Id})$ for some $k_1, k_2 \in \mathbb{N}$. Then, by Lemma 3.4, $g = (x^i y^j)^{2^{k_2}} x^{s2^{k_1} \alpha^{2^{k_2} j}}$. Now, suppose that $\alpha \in \mathfrak{C}_1$, then $\alpha = 2^\beta$ for some $\beta \in \mathbb{N}$. Then, $g = (x^i y^j)^{2^{k_2}} x^{s2^{k_1} 2^{\beta 2^{k_2} j}} = (x^i y^j)^{2^{k_2}} x^{s2^{k_1 + \beta 2^{k_2} j}} \in \text{supp}(\langle x^i y^j \rangle_{Id} \langle x^s \rangle_{Id})$. Hence, $\text{supp}(\langle x^s \rangle_{Id} \langle x^i y^j \rangle_{Id}) \subseteq \text{supp}(\langle x^i y^j \rangle_{Id} \langle x^s \rangle_{Id})$. By a similar argument, it can be verified that $\text{supp}(\langle x^i y^j \rangle_{Id} \langle x^s \rangle_{Id}) \subseteq \text{supp}(\langle x^s \rangle_{Id} \langle x^i y^j \rangle_{Id})$ and thus we have that $\langle x^s \rangle_{Id} \langle x^i y^j \rangle_{Id} = \langle x^i y^j \rangle_{Id} \langle x^s \rangle_{Id}$. Therefore, $\langle x^s \rangle_{Id} \in J(G)$. \square

As a whole, for non-abelian $G = C_p \rtimes C_q = \langle x, y | x^p = y^q = 1, xy = yx^\alpha \rangle$, if $\alpha \in \mathfrak{C}_1$, then $J(G) = \{\langle x^i y^j \rangle_{Id} | j = 0\}$.

3.2 Dihedral Groups

Recall the dihedral groups $D_{2n} = C_n \rtimes C_2 = \langle a, b | a^n = b^2 = 1, ab = ba^{-1} \rangle$ with $|D_{2n}| = 2n$. Every $g \in D_{2n}$ is either in the form of a^k or $a^k b$ for some $k \in \mathbb{N}$. For $g = a^k b$, note that $\text{ord}(g) = 2$.

Theorem 3.6 Let $D_{2n} = \langle a, b | a^n = b^2 = 1, ab = ba^{-1} \rangle$. Then, $I_{\langle Id \rangle}(D_{2n}) = I_{\langle Id \rangle}(C_{\langle a \rangle})$.

Proof For each $k \in \mathbb{N}$, as $\text{ord}(a^k b) = 2$, $a^k b$ cannot be a generator for any $\langle g \rangle_{Id} \in I_{\langle Id \rangle}(D_{2n})$ by Theorem 1.5. This results in $\langle g \rangle_{Id} \in I_{\langle Id \rangle}(D_{2n})$ if and only if $g \in C_{\langle a \rangle}$ with $\text{ord}(g)$ odd. Thus, it follows from Theorem 1.5 that $\langle g \rangle_{Id} \in I_{\langle Id \rangle}(D_{2n})$ if and only if $I_{\langle Id \rangle}(C_{\langle a \rangle})$. \square

The following corollary is a result from Theorem 1.4 and Theorem 3.6.

Corollary 3.7 Let $D_{2n} = \langle a, b | a^n = b^2 = 1, ab = ba^{-1} \rangle$. Then, $J(D_{2n}) = I_{\langle Id \rangle}(D_{2n})$.

We provide an example to illustrate the case when $n = 7$.

Example 3.8 Let $D_{14} = \langle a, b | a^7 = b^2 = 1, ab = ba^{-1} \rangle$. It can be verified that there are three generated idempotents in $F_2 C_{\langle a \rangle}$, namely 1 , $\langle a \rangle_{Id} = a + a^2 + a^4$ and $\langle a^3 \rangle_{Id} = a^3 + a^5 + a^6$. Thus, $I_{\langle Id \rangle}(C_{\langle a \rangle}) = \{1, \langle a \rangle_{Id}, \langle a^3 \rangle_{Id}\}$. By Theorem 3.6, $I_{\langle Id \rangle}(D_{14}) = I_{\langle Id \rangle}(C_{\langle a \rangle})$. Then, it follows from Corollary 3.7 that $J(D_{14}) = I_{\langle Id \rangle}(D_{14}) = \{1, \langle a \rangle_{Id}, \langle a^3 \rangle_{Id}\}$.

3.3 Non-abelian Groups of Order $2^k p$

Let G be of order $2^k p$, where $k \in \mathbb{Z}^+$ and $p > 2$ is a prime. Let α be the total number of Sylow p -subgroups of G . The result on $J(G)$ depends on the uniqueness of the Sylow p -subgroup of G .

Proposition 3.9 *Let $|G| = 2^k p$ where p is a prime such that $p > 2$.*

1. *If the Sylow p -subgroup of G is unique, then $J(G) = I_{\langle Id \rangle}(G)$.*
2. *If the Sylow p -subgroups of G are not unique, then $J(G) = \{1\}$.*

Proof Suppose that C_p is the unique Sylow p -subgroup of G , note that C_p is the only odd order subgroup of G . By Theorem 1.5, all $\langle g \rangle_{Id} \in I_{\langle Id \rangle}(G)$ have generator $g \in C_p$. Then, $I_{\langle Id \rangle}(G) = I_{\langle Id \rangle}(C_p)$. By Theorem 1.4, $J(G) = I_{\langle Id \rangle}(G)$.

On the other hand, suppose that the Sylow p -subgroups of G are not unique. Let $\langle g \rangle_{Id} \in I_{\langle Id \rangle}(G) \setminus \{1\}$. Then, $ord(g)$ must be odd by Theorem 1.5. Hence, $g \in C_p$ for some Sylow p -subgroup C_p of G . Note that $C_{\langle g \rangle} = C_p$ is not a normal subgroup of G . Also, $C_{\langle g \rangle}$ is a maximal odd order subgroup of G . Thus, $C_{\langle g \rangle}$ is not an odd-subnormal subgroup of G . Then, it follows from Theorem 2.9 that $\langle g \rangle_{Id} \notin J(G)$. Hence, $J(G) = \{1\}$. \square

This characterization of $J(G)$ leads to the corollary for the case where $|G| = 2^k p$ and p is a prime such that $p > 2^k$.

Corollary 3.10 *Let $|G| = 2^k p$ with $p > 2^k$ is a prime. Then, $J(G) = I_{\langle Id \rangle}(G)$.*

Proof We claim that the Sylow p -subgroup of G is unique and thus the result follows from Proposition 3.9. Suppose that there exist P_1 and P_2 , distinct Sylow p -subgroups of G , note that $P_1 \cap P_2 = \{1\}$. Then, $P_1 P_2$ has p^2 distinct elements. This results in $|G| > p^2 > 2^k p$. \square

4 Equivalency of Zero-divisor Codes in $F_2 G$

Throughout this section, G denotes a finite group (not necessary non-abelian). Results in Section 2,3 and [9] imply that most of the groups G have a pool of non-trivial idempotents to serve as the generator for their respective zero-divisor codes in $F_2 G$. This section gives a comprehensive study on the equivalency of general zero-divisor codes in $F_2 G$ based on Definition 1.2, where $W_1 u_1$ and $W_2 u_2$ are said to be equivalent if there exists an F_2 -linear isomorphism $\varphi : W_1 u_1 \rightarrow W_2 u_2$, such that for every $c \in W_1 u_1$, $wt(c) = wt((c)\varphi)$.

The property $wt(c) = wt((c)\varphi)$ for every $c \in W_1 u_1$ of φ is called weight-preserving. For every Wu , the notation $supp(Wu)$ denotes $\bigcup_{c \in Wu} supp(c)$. For the remaining paper, $W_1 u_1$ and $W_2 u_2$ are zero-divisor codes in $F_2 G_1$ and $F_2 G_2$ respectively, with $supp(W_1 u_1) \subseteq G_1$ and $supp(W_2 u_2) \subseteq G_2$.

We first show that the existence of an injective function $\chi : supp(W_1 u_1) \rightarrow supp(W_2 u_2)$ associated to a function $\varphi : W_1 u_1 \rightarrow W_2 u_2$ in such a way that $(\sum_{g \in supp(W_1 u_1)} a_g g)\varphi = \sum_{g \in supp(W_1 u_1)} a_g \chi(g)$ plays a vital role in making φ a weight-preserving injective F_2 -linear map.

Theorem 4.1 *Let W_1u_1 and W_2u_2 be zero-divisor codes in F_2G_1 and F_2G_2 respectively. If there exists a function $\varphi : W_1u_1 \rightarrow W_2u_2$ in such a way that $(\sum_{g \in \text{supp}(W_1u_1)} a_g g)\varphi = \sum_{g \in \text{supp}(W_1u_1)} a_g \chi(g)$ for some function $\chi : \text{supp}(W_1u_1) \rightarrow \text{supp}(W_2u_2)$, then φ is a weight-preserving injective F_2 -linear map if and only if χ is injective.*

Proof Suppose that χ is injective. Then, for each $c = \sum_{g \in \text{supp}(W_1u_1)} a_g g \in W_1u_1$,

$$wt((c)\varphi) = wt\left(\sum_{g \in \text{supp}(W_1u_1)} a_g \chi(g)\right) = wt\left(\sum_{g \in \text{supp}(W_1u_1)} a_g g\right) = wt(c).$$

Thus, φ is a weight-preserving function. Note that φ is closed under scalar multiplication over F_2 and for $c_1 = \sum_{g \in \text{supp}(W_1u_1)} a_g g, c_2 = \sum_{g \in \text{supp}(W_1u_1)} b_g g \in W_1u_1$,

$$\begin{aligned} (c_1 + c_2)\varphi &= \sum_{g \in \text{supp}(W_1u_1)} (a_g + b_g)\chi(g) \\ &= \sum_{g \in \text{supp}(W_1u_1)} a_g \chi(g) + \sum_{g \in \text{supp}(W_1u_1)} b_g \chi(g) \\ &= (c_1)\varphi + (c_2)\varphi \end{aligned}$$

Hence, φ is an F_2 -linear map with:

$$\text{Ker}(\varphi) = \left\{ \sum_{g \in \text{supp}(W_1u_1)} a_g g \in W_1u_1 \mid \left(\sum_{g \in \text{supp}(W_1u_1)} a_g g\right)\varphi = 0 \right\}.$$

By the choice of φ , note that $\text{Ker}(\varphi) = \{0\}$, thus φ is injective.

Conversely, suppose that χ is not injective, there exist $g_1, g_2 \in \text{supp}(W_1u_1)$, $g_1 \neq g_2$ such that $\chi(g_1) = \chi(g_2)$. Note that as W_1u_1 is a linear code, there always exists $v \in W_1u_1$ such that $g_1, g_2 \in \text{supp}(v)$. Then, $wt((v)\varphi) \leq wt(v) - 1$, thus φ is not weight-preserving. \square

The following corollary is a direct result from Theorem 4.1.

Corollary 4.2 *Let W_1u_1 and W_2u_2 be zero-divisor codes in F_2G_1 and F_2G_2 respectively. If there exists a function $\varphi : W_1u_1 \rightarrow W_2u_2$ in such a way that $(\sum_{g \in \text{supp}(W_1u_1)} a_g g)\varphi = \sum_{g \in \text{supp}(W_1u_1)} a_g \chi(g)$ for some function $\chi : \text{supp}(W_1u_1) \rightarrow \text{supp}(W_2u_2)$, then W_1u_1 is equivalent to $\text{Im}(\varphi)$ if and only if χ is injective.*

Consider $W_1u_1 = F_2G_1u_1$ and $W_2u_2 = F_2G_2u_2$ from the family of ideal zero-divisor codes in F_2G_1 and F_2G_2 respectively. Then, $\text{supp}(F_2G_1u_1) = G_1$ and $\text{supp}(F_2G_2u_2) = G_2$. Let $\chi : G_1 \rightarrow G_2$ be an injective function. To utilize Corollary 4.2, the focus is next concentrated on the sufficient conditions for the existence of φ in the form of:

$$\left(\sum_{g \in G_1} a_g g\right)\varphi = \sum_{g \in G_1} a_g \chi(g) \quad (4.1)$$

Let $c = (\sum_{g \in G_1} a_g g)u_1 \in F_2G_1u_1$. Note that $c = \sum_{g \in G_1} \sum_{g_1 \in \text{supp}(u_1)} a_g g g_1 = \sum_{g \in G_1, a_g \neq 0} \sum_{g_1 \in \text{supp}(u_1)} g g_1$. Then, we have:

$$(c)\varphi = \left(\sum_{g \in G_1, a_g \neq 0} \sum_{g_1 \in \text{supp}(u_1)} g g_1 \right) \varphi = \sum_{g \in G_1, a_g \neq 0} \sum_{g_1 \in \text{supp}(u_1)} \chi(g g_1) \quad (4.2)$$

On the other hand, in order to have $(c)\varphi \in F_2G_2u_2$:

$$(c)\varphi = \left(\sum_{h \in G_2} b_h h \right) u_2 = \sum_{h \in G_2} \sum_{g_2 \in \text{supp}(u_2)} b_h h g_2 = \sum_{h \in G_2, b_h \neq 0} \sum_{g_2 \in \text{supp}(u_2)} h g_2. \quad (4.3)$$

Then, by (4.2) and (4.3),

$$\sum_{g \in G_1, a_g \neq 0} \sum_{g_1 \in \text{supp}(u_1)} \chi(g g_1) = \sum_{h \in G_2, b_h \neq 0} \sum_{g_2 \in \text{supp}(u_2)} h g_2 \quad (4.4)$$

In addition, suppose that χ further satisfies the following conditions: For every $g \in G_1$ and $g_1 \in \text{supp}(u_1)$, $\chi(g g_1) = \chi(g)\chi(g_1)$ and $\chi(\text{supp}(u_1)) = \text{supp}(u_2)$. Then, from (4.4):

$$\begin{aligned} (c)\varphi &= \sum_{g \in G_1, a_g \neq 0} \sum_{g_1 \in \text{supp}(u_1)} \chi(g g_1) \\ &= \sum_{g \in G_1, a_g \neq 0} \sum_{g_1 \in \text{supp}(u_1)} \chi(g)\chi(g_1) \\ &= \sum_{g \in G_1, a_g \neq 0} \chi(g) \sum_{g_1 \in \text{supp}(u_1)} \chi(g_1) \\ &= \left(\sum_{g \in G_1, a_g \neq 0} \chi(g) \right) u_2 \end{aligned}$$

These two conditions on χ guarantee the existence of $\varphi : F_2G_1u_1 \rightarrow F_2G_2u_2$ in such a way that $(\sum_{g \in G_1} a_g g)\varphi = \sum_{g \in G_1} a_g \chi(g)$ for a given injective function $\chi : G_1 \rightarrow G_2$. Furthermore, if χ is surjective, notice that $(c)\varphi = (\sum_{g \in G_1, a_g \neq 0} \chi(g))u_2$ also indicates that φ is onto. A bijective χ that satisfies each of the two properties imposed earlier is defined formally.

Definition 4.3 Let u_1 and u_2 be zero-divisors in F_2G_1 and F_2G_2 respectively. Let $\chi : A \rightarrow B$ be a bijective function for $A \subseteq G_1$ and $B \subseteq G_2$. Then:

1. χ is denoted as χ_{u_1, u_2} if $\chi(\text{supp}(u_1)) = \text{supp}(u_2)$.
2. χ is termed a u_1 -homomorphism if $\chi(gh) = \chi(g)\chi(h)$ for every $g \in A$ and $h \in \text{supp}(u_1)$.

Using the terms in Definition 4.3, Theorem 4.1 and Corollary 4.2, the following two results naturally hold.

Theorem 4.4 *Let u_1 and u_2 be zero-divisors in F_2G_1 and F_2G_2 respectively. If $\chi_{u_1, u_2} : G_1 \rightarrow G_2$ is a u_1 -homomorphism, then there exists a weight-preserving F_2 -linear isomorphism $\varphi : F_2G_1u_1 \rightarrow F_2G_2u_2$ in such a way that $(\sum_{g \in G_1} a_g g)\varphi =$*

$$\sum_{g \in G_1} a_g \chi_{u_1, u_2}(g).$$

Corollary 4.5 *Let u_1 and u_2 be zero-divisors in F_2G_1 and F_2G_2 respectively. If there exists $\chi_{u_1, u_2} : G_1 \rightarrow G_2$ which is a u_1 -homomorphism, then $F_2G_1u_1$ and $F_2G_2u_2$ are equivalent ideal zero-divisor codes.*

For each F_2 -submodule W_1 of F_2G_1 , W_1u_1 is a subcode of $F_2G_1u_1$. Suppose that $F_2G_1u_1$ and $F_2G_2u_2$ are equivalent by a weight-preserving F_2 -linear isomorphism $\varphi : F_2G_1u_1 \rightarrow F_2G_2u_2$. Define φ_\downarrow by restricting φ to a smaller domain W_1u_1 , and then it is easy to see that the equivalency of $F_2G_1u_1$ and $F_2G_2u_2$ implies the equivalency of W_1u_1 and $Im(\varphi_\downarrow)$, a subcode of $F_2G_2u_2$. The following result gives the detail of $Im(\varphi_\downarrow)$ when $(\sum_{g \in G_1} a_g g)\varphi = \sum_{g \in G_1} a_g \chi_{u_1, u_2}(g)$ with respect to a u_1 -homomorphism $\chi_{u_1, u_2} : G_1 \rightarrow G_2$.

Corollary 4.6 *Let u_1 and u_2 be zero-divisors in F_2G_1 and F_2G_2 respectively. Let W_1 be an F_2 -submodule of F_2G_1 with basis α such that αu_1 is a basis for W_1u_1 . If there exists a u_1 -homomorphism $\chi_{u_1, u_2} : G_1 \rightarrow G_2$, then W_1u_1 is equivalent to W_2u_2 , where W_2 is an F_2 -submodule of F_2G_2 with a basis in the form of $\beta = \{ \sum_{g \in \text{supp}(v)} \chi_{u_1, u_2}(g) | v \in \alpha \}$.*

Proof Note that $F_2G_1u_1$ and $F_2G_2u_2$ are equivalent with respect to some weight-preserving F_2 -linear isomorphism $\varphi : F_2G_1u_1 \rightarrow F_2G_2u_2$ in such a way that $((\sum_{g \in G_1} a_g g)u_1)\varphi = (\sum_{g \in G_1} a_g \chi_{u_1, u_2}(g))u_2$ for u_1 -homomorphism $\chi_{u_1, u_2} : G_1 \rightarrow G_2$.

Define φ_\downarrow to be the restriction of φ onto the smaller domain W_1u_1 and also $\chi_{u_1, u_2}' : \text{supp}(W_1u_1) \rightarrow \text{supp}(Im(\varphi_\downarrow))$ by restricting the domain and codomain of χ_{u_1, u_2} to $\text{supp}(W_1u_1) \subseteq G_1$ and $\text{supp}(Im(\varphi_\downarrow)) \subseteq G_2$. Clearly, φ_\downarrow is a weight preserving injective F_2 -linear map such that

$$(\sum_{g \in \text{supp}(W_1u_1)} a_g g)\varphi_\downarrow = \sum_{g \in \text{supp}(W_1u_1)} a_g \chi_{u_1, u_2}'(g)$$

and χ_{u_1, u_2}' is a u_1 -homomorphism. Hence, W_1u_1 is a code with basis αu_1 , with $|\alpha u_1| = |\alpha|$ and is equivalent to $Im(\varphi_\downarrow)$. Note that $(c)\varphi_\downarrow = ((\sum_{v \in \alpha} a_v v)u_1)\varphi_\downarrow =$

$$((\sum_{v \in \alpha} a_v \sum_{g \in \text{supp}(v)} g)u_1)\varphi_\downarrow = (\sum_{v \in \alpha} a_v \sum_{g \in \text{supp}(v)} \chi_{u_1, u_2}'(g))u_2.$$

Let $W_2u_2 = Im(\varphi_\downarrow)$ where W_2 is an F_2 -submodule of F_2G_2 . Clearly, W_2 is spanned by $\beta = \{ \sum_{g \in \text{supp}(v)} \chi_{u_1, u_2}'(g) | v \in \alpha \}$. Then, βu_2 contains a basis for W_2u_2 .

However, note that $|\beta u_2| = |\beta| = |\alpha| = |\alpha u_1|$ as χ_{u_1, u_2}' is bijective. As φ_\downarrow is an isomorphism, $\gamma = \beta u_2$ is linearly independent over F_2 and thus γ is a basis for $Im(\varphi_\downarrow) = W_2u_2$. Hence, β is a basis for W_2u_2 otherwise $|\gamma| > |\beta|$. \square

5 Equivalency of Zero-divisor Codes in F_2G with Idempotent Generator

The studies of zero-divisor codes with idempotent generators are discussed from this point onwards. Classically, cyclic and dihedral groups with the same order $2n$

for $n > 2$ can never be isomorphic as commutativity holds for the former but not latter. In the next subsection, we answer the earlier conjecture partially, that is having our classified idempotent as generator, each zero-divisor code in F_2D_{2n} is equivalent to some zero-divisor code in F_2C_{2n} .

5.1 Zero-divisor Codes in F_2D_{2n} and F_2C_{2n}

Throughout this subsection, we denote $C_{2n} = C_{\langle x \rangle}$ as a cyclic group of order $2n$ and $D_{2n} = \langle a, b | a^n = b^2 = 1, ab = ba^{-1} \rangle$ as a dihedral group of order $2n$. Let $n = 2^\alpha \beta$ for some integer $\alpha \geq 0$ and odd integer $\beta > 1$, then $I_{\langle Id \rangle}(C_{2n}) = I_{\langle Id \rangle}(C_{\langle x^{2^{2^\alpha}} \rangle})$ by Theorem 1.5. On the other hand, $I_{\langle Id \rangle}(D_{2n}) = I_{\langle Id \rangle}(C_{\langle a \rangle}) = I_{\langle Id \rangle}(C_{\langle a^{2^\alpha} \rangle})$ by Theorem 3.6. We thus obtain a one-to-one mapping from $I(C_{2n})$ to $I(D_{2n})$ as follows.

Proposition 5.1 *Let $I_{\langle Id \rangle}(C_{2n}) = \{\langle x^{2^{s_i}} \rangle_{Id} | i \in \{1, 2, \dots, l\}\}$ and $\alpha_1, \alpha_2, \dots, \alpha_l \in F_2$. Define $\rho : I(C_{2n}) \rightarrow I(D_{2n})$ by $\rho(\sum_{i=1}^l \alpha_i \langle x^{2^{s_i}} \rangle_{Id}) = \sum_{i=1}^l \alpha_i \langle a^{s_i} \rangle_{Id}$. Then, ρ is a one-to-one mapping.*

Setting each pair of idempotents in Proposition 5.1 as generator of ideal zero-divisor codes in F_2C_{2n} and F_2D_{2n} respectively, the two resultant codes are always equivalent as shown by the following theorem.

Theorem 5.2 *Let $I_{\langle Id \rangle}(C_{2n}) = \{\langle x^{2^{s_i}} \rangle_{Id} | i \in \{1, 2, \dots, l\}\}$ and $\alpha_1, \alpha_2, \dots, \alpha_l \in F_2$. Consider $e_1 = \sum_{i=1}^l \alpha_i \langle a^{s_i} \rangle_{Id} \in I(D_{2n})$ and $e_2 = \sum_{i=1}^l \alpha_i \langle x^{2^{s_i}} \rangle_{Id} \in I(C_{2n})$. Then, $F_2D_{2n}e_1$ and $F_2C_{2n}e_2$ are equivalent ideal zero-divisor codes.*

Proof First, note that $\{(2i+1) \bmod 2n | i \in \{0, 1, \dots, n-1\}\} = \{(-2i-1) \bmod 2n | i \in \{0, 1, \dots, n-1\}\}$ and $(x^{2^{i+1}})^{-1} = x^{-2^{i-1}}$. Then, C_{2n} can be partitioned as follows

$$\begin{aligned} C_{2n} &= \{x^{2^i} | i \in \{0, 1, \dots, n-1\}\} \cup \{x^{2^{i+1}} | i \in \{0, 1, \dots, n-1\}\} \\ &= \{x^{2^i} | i \in \{0, 1, \dots, n-1\}\} \cup \{x^{-2^{i-1}} | i \in \{0, 1, \dots, n-1\}\}. \end{aligned}$$

Define $\chi : D_{2n} \rightarrow C_{2n}$ such that

$$\chi(a^i b^j) = \begin{cases} x^{-2^{i-1}} & j = 1 \\ x^{2^i} & j = 0 \end{cases} \quad (5.1)$$

Clearly, χ is bijective and $\chi(\text{supp}(e_1)) = \text{supp}(e_2)$. Hence, $\chi = \chi_{e_1, e_2}$. Let $a^{2^k s_j} \in \text{supp}(e_1)$ for some $k \in \mathbb{Z}^+$. Then, for every $a^i b \in D_{2n}$, $\chi((a^i b)(a^{2^k s_j}))$ can be written as:

$$\chi(a^{i-2^k s_j} b) = x^{-2(i-2^k s_j)-1} = x^{-2i-1} x^{2(2^k s_j)} = \chi(a^i b) \chi(a^{2^k s_j}).$$

Similarly, for every $a^i \in D_{2n}$:

$$\chi(a^i a^{2^k s_j}) = \chi(a^{i+2^k s_j}) = x^{2(i+2^k s_j)} = x^{2i} x^{2(2^k s_j)} = \chi(a^i) \chi(a^{2^k s_j}).$$

This concludes that χ_{e_1, e_2} is an e_1 -homomorphism, thus the proof follows from Corollary 4.5. \square

Example 5.3 Utilizing Proposition 5.1, we establish a one-to-one correspondence from the idempotents in Example 1.6 to Example 3.8, a listing of all ideal zero-divisor codes in F_2C_{14} with idempotent generator and its equivalent partner in F_2D_{14} is obtained.

Ideal Zero-divisor Codes in F_2C_{14}	Ideal Zero-divisor Codes in F_2D_{14}
F_2C_{14}	F_2D_{14}
$F_2C_{14}\langle x^2 \rangle_{Id}$	$F_2D_{14}\langle a \rangle_{Id}$
$F_2C_{14}\langle x^6 \rangle_{Id}$	$F_2D_{14}\langle a^3 \rangle_{Id}$
$F_2C_{14}(1 + \langle x^2 \rangle_{Id})$	$F_2D_{14}(1 + \langle a \rangle_{Id})$
$F_2C_{14}(1 + \langle x^6 \rangle_{Id})$	$F_2D_{14}(1 + \langle a^3 \rangle_{Id})$
$F_2C_{14}(\langle x^2 \rangle_{Id} + \langle x^6 \rangle_{Id})$	$F_2D_{14}(\langle a \rangle_{Id} + \langle a^3 \rangle_{Id})$
$F_2C_{14}(1 + \langle x^2 \rangle_{Id} + \langle x^6 \rangle_{Id})$	$F_2D_{14}(1 + \langle a \rangle_{Id} + \langle a^3 \rangle_{Id})$

Therefore, it follows that for each zero-divisor code in F_2C_{2n} in the form of W_1e_1 , the equivalent partner in F_2D_{2n} is W_2e_2 , where W_2 can be obtained by Corollary 4.6.

5.2 Zero-divisor Codes in F_2C_n

In this last subsection, we continue the study of the equivalency of zero-divisor codes in F_2C_n with idempotent generator. Note that for every $\langle x^s \rangle_{Id} \in I_{(Id)}(C_n)$, $ord(x^s)$ must be odd by Theorem 1.5. Consider distinct idempotents $e_1 = \langle x^{s_1} \rangle_{Id}, e_2 = \langle x^{s_2} \rangle_{Id} \in I_{(Id)}(C_n)$. For the case where n is odd and $ord(x^{s_1}) = ord(x^{s_2}) = n$, both x^{s_1} and x^{s_2} are generators of C_n and thus heuristically an e_1 -homomorphism $\chi_{e_1, e_2} : C_n \rightarrow C_n$ is the automorphism of C_n defined by $\chi_{e_1, e_2}(x^{s_1^i}) = x^{s_2^i}$. Then, by Corollary 4.5, the next result holds.

Proposition 5.4 Let $x^{s_1}, x^{s_2} \in C_n$ for odd n . If $ord(x^{s_1}) = ord(x^{s_2}) = n$, then $F_2C_n\langle x^{s_1} \rangle_{Id}$ and $F_2C_n\langle x^{s_2} \rangle_{Id}$ are equivalent ideal zero-divisor codes in F_2C_n .

We summarize the case where n is an odd prime in the next corollary.

Corollary 5.5 Let n be an odd prime. All $F_2C_n\langle x^s \rangle_{Id}$ with $s \neq 0$ are equivalent ideal zero-divisor codes in F_2C_n .

The establishment of equivalency between $F_2C_n\langle x^s \rangle_{Id}$ depends ultimately on the property of $\langle x^s \rangle_{Id}$. Gaining insight from Proposition 5.4, we prove that $ord(x^s)$ categorizes equivalent $F_2C_n\langle x^s \rangle_{Id}$. Before that, the following lemma is needed.

Lemma 5.6 For distinct $\langle x^{s_1} \rangle_{Id}, \langle x^{s_2} \rangle_{Id} \in I_{(Id)}(C_n)$, if $ord(x^{s_1}) = ord(x^{s_2})$, then $|supp\langle x^{s_1} \rangle_{Id}| = |supp\langle x^{s_2} \rangle_{Id}|$.

Proof Note that for $i \in \{1, 2\}$, $|supp\langle x^{s_i} \rangle_{Id}|$ is the smallest non-negative integer k_i such that $ord(x^{s_i}) | 2^{k_i} - 1$. Then, it follows from $ord(x^{s_1}) = ord(x^{s_2})$ that $k_1 = k_2$. \square

Theorem 5.7 For distinct $e_1 = \langle x^{s_1} \rangle_{Id}, e_2 = \langle x^{s_2} \rangle_{Id} \in I_{(Id)}(C_n)$, $ord(x^{s_1}) = ord(x^{s_2})$ if and only if there exists an e_1 -homomorphism $\chi_{e_1, e_2} : C_n \rightarrow C_n$.

Proof Suppose that $\text{ord}(x^{s_1}) = \text{ord}(x^{s_2})$. Then, $|\text{supp}\langle x^{s_1} \rangle_{Id}| = |\text{supp}\langle x^{s_2} \rangle_{Id}|$ by Lemma 5.6. Let $\text{ord}(x^{s_1}) = \text{ord}(x^{s_2}) = m$, then by Lagrange's Theorem, $mr = n$ for some $r \in \mathbb{Z}^+$. Thus, $C_n = \{x^{i+j s_1} \mid i \in \{0, 1, \dots, r-1\}, j \in \{0, 1, \dots, m-1\}\}$. Hence, $\chi_{e_1, e_2} : C_n \rightarrow C_n$ can be defined as $\chi_{e_1, e_2}(x^{i+j s_1}) = x^{i+j s_2}$ for each $i \in \{0, 1, \dots, r-1\}$ and $j \in \{0, 1, \dots, m-1\}$. Note that χ_{e_1, e_2} is also an e_1 -homomorphism as for every $t = i + j s_1 \in \{0, 1, \dots, n-1\}$ and $k \in \mathbb{Z}^+$, $\chi_{e_1, e_2}(x^t x^{s_1 2^k}) = \chi_{e_1, e_2}(x^{i+s_1(j+2^k)})$. We claim that $\chi_{e_1, e_2}(x^{i+s_1(j+2^k)}) = x^{i+s_2(j+2^k)}$. To prove the claim, it is sufficient to show that for $h \in \{0, 1, \dots, m-1\}$ and $h' \in \mathbb{N} \setminus \{0, 1, \dots, m-1\}$, $x^{i+h s_2} = x^{i+h' s_2}$ if $x^{i+h s_1} = x^{i+h' s_1}$. Let $h' = h + \alpha$, then:

$$\begin{aligned} x^{i+h s_1} &= x^{i+h' s_1} \\ x^{h s_1} &= x^{h' s_1} \\ x^{h s_1} &= x^{(h+\alpha) s_1} \\ 1 &= x^{\alpha s_1} \end{aligned}$$

This implies that α must be a multiple of m and thus:

$$\begin{aligned} x^{i+h' s_2} &= x^{i+(h+\alpha) s_2} \\ &= x^{i+h s_2} x^{\alpha s_2} \\ &= x^{i+h s_2} (1) \\ &= x^{i+h s_2}. \end{aligned}$$

This results in:

$$\begin{aligned} \chi_{e_1, e_2}(x^{i+s_1(j+2^k)}) &= x^{i+s_2(j+2^k)} \\ &= x^{i+j s_2} x^{s_2 2^k} \\ &= \chi_{e_1, e_2}(x^t) \chi_{e_1, e_2}(x^{s_1 2^k}). \end{aligned}$$

Conversely, suppose that there exists an e_1 -homomorphism $\chi_{e_1, e_2} : C_n \rightarrow C_n$. Without loss of generality, assume that $m_1 = \text{ord}(x^{s_1}) < \text{ord}(x^{s_2}) = m_2$ where both m_1 and m_2 are odd. Note that $\chi_{e_1, e_2}(x^{s_1}) = x^{2^k s_2}$ for some $k \in \mathbb{N}$. It can be shown by mathematical induction that $\chi_{e_1, e_2}((x^{s_1})^t) = (\chi_{e_1, e_2}(x^{s_1}))^t$ for every $t \in \mathbb{N}$. Hence,

$$\chi_{e_1, e_2}(1) = \chi_{e_1, e_2}((x^{s_1})^{m_1}) = (\chi_{e_1, e_2}(x^{s_1}))^{m_1} = (x^{2^k s_2})^{m_1} \neq 1$$

as $\text{ord}(x^{2^k s_2}) = \text{ord}(x^{s_2}) = m_2 > m_1$. This is a contradiction to $\chi_{e_1, e_2}(x^{s_1}) = x^{2^k s_2}$ for some $k \in \mathbb{N}$ as

$$\chi_{e_1, e_2}(x^{s_1}) = \chi_{e_1, e_2}(1 x^{s_1}) = \chi_{e_1, e_2}(1) \chi_{e_1, e_2}(x^{s_1}) = \chi_{e_1, e_2}(1) x^{2^k s_2} \neq x^{2^k s_2}.$$

Thus, $\text{ord}(x^{s_1}) = \text{ord}(x^{s_2})$. \square

Together with Corollary 4.5, Theorem 5.7 implies the following result.

Corollary 5.8 *For distinct $e_1 = \langle x^{s_1} \rangle_{Id}$, $e_2 = \langle x^{s_2} \rangle_{Id} \in I_{(Id)}(C_n)$, if $\text{ord}(x^{s_1}) = \text{ord}(x^{s_2})$, then $F_2 C_n e_1$ and $F_2 C_n e_2$ are equivalent ideal zero-divisor codes in $F_2 C_n$.*

For zero-divisor codes in F_2C_n having a finite sum of some generated idempotents as generator, it is not easy to define an e_1 -homomorphism $\chi_{e_1, e_2} : C_n \rightarrow C_n$ explicitly for each n . We end this section with the following subcase.

Proposition 5.9 *Let $l = |I_{(Id)}(C_n)|$ and $\alpha_1, \alpha_2, \dots, \alpha_l \in F_2$. Consider distinct $e_1 = \sum_{i=1}^l \alpha_i \langle x^{s_i} \rangle_{Id}, e_2 = \sum_{i=1}^l \alpha_i \langle x^{s_i k} \rangle_{Id} \in I(C_n)$ with $(n, k) = 1$. Then, $F_2C_n e_1$ and $F_2C_n e_2$ are equivalent ideal zero-divisor codes in F_2C_n .*

Proof As $(n, k) = 1$, define an e_1 -homomorphism $\chi : C_n \rightarrow C_n$ as the automorphism that maps the generator x to the generator x^k , that is $\chi(x^i) = x^{ik}$. Note that $\chi(\text{supp}(e_1)) = \text{supp}(e_2)$. Thus, $\chi = \chi_{e_1, e_2}$ and Corollary 4.5 affirms the equivalency between $F_2C_n e_1$ and $F_2C_n e_2$. \square

6 Conclusion

This paper is on the usefulness of idempotent generators in establishing the equivalency of zero-divisor codes in F_2G , via their classifications on $I(G)$. For non-abelian G , the fact that $I(G)$ need not form a vector space over F_2 triggered the study of the *joinability* between distinct $\langle g \rangle_{Id}, \langle h \rangle_{Id} \in I_{(Id)}(G)$. The study of $I_{(Id)}(G)$ was mainly focused on its special subset $J(G) = \{\langle g \rangle_{Id} \in I_{(Id)}(G) \mid \langle g \rangle_{Id} + \langle h \rangle_{Id} \in I(G)\}$. The condition that $C_{\langle g \rangle}$ permutes with $C_{\langle h \rangle}$ was proved to be necessary for $\langle g \rangle_{Id} + \langle h \rangle_{Id} \in I(G)$, and so for $\langle g \rangle_{Id} \in J(G)$, $C_{\langle g \rangle}$ must be odd-permutable and thus odd-subnormal in G . These results were then used to study $J(G)$ explicitly for groups of order pq , $2^k p$ and dihedral groups.

The non-triviality of $J(G)$ for these several families of G ensures that there is a large pool of idempotents to serve as the generator of zero-divisor codes in F_2G in order to study the equivalency between the resultant codes. In search of the weight-preserving F_2 -linear isomorphism $\varphi : F_2G_1 u_1 \rightarrow F_2G_2 u_2$, we instinctively concentrated on those which are expressible in terms of $\chi : G_1 \rightarrow G_2$, that is $(\sum_{g \in G_1} a_g g) \varphi = \sum_{g \in G_1} a_g \chi(g)$. The exhibition of φ being weight-preserving F_2 -linear isomorphic is then completely dependent on the properties of χ . Precisely, the existence of a bijective χ satisfying $\chi(\text{supp}(u_1)) = \text{supp}(u_2)$ simultaneously being a u_1 -homomorphism suffices to guarantee that φ a weight-preserving F_2 -linear isomorphism.

The existence of a weight-preserving F_2 -linear isomorphism φ can be easily shown when both the generators u_1 and u_2 are our classified idempotents. Our previous classification of $I(C_{2n})$ and $I(D_{2n})$ amalgamated into the one-to-one mapping: $\sum_{i=1}^l \alpha_i \langle x^{2s_i} \rangle_{Id} \in I(C_{2n})$ if and only if $\sum_{i=1}^l \alpha_i \langle a^{s_i} \rangle_{Id} \in I(D_{2n})$. Showing the existence of χ in (5.1), we then prove that $F_2C_{2n}(\sum_{i=1}^l \alpha_i \langle x^{2s_i} \rangle_{Id})$ and $F_2D_{2n}(\sum_{i=1}^l \alpha_i \langle a^{s_i} \rangle_{Id})$ always form a pair of equivalent ideal zero-divisor codes. Hence, by Corollary 4.6, we show the conjecture ‘‘Every group ring code in F_2D_{2n} is equivalent to some group ring codes in F_2C_{2n} ’’ in the cases where the code generators are idempotents in the form illustrated in Theorem 1.7.

Last but not least, the order of the generator contributed to an elegant categorization on the equivalency of $F_2C_n\langle x^s \rangle_{Id}$. If $\langle x^{s_1} \rangle_{Id}, \langle x^{s_2} \rangle_{Id} \in I_{\langle Id \rangle}(C_n)$ such that $ord(x^{s_1}) = ord(x^{s_2})$, then the desirable χ exists, resulting in $F_2C_n\langle x^{s_1} \rangle_{Id}$ and $F_2C_n\langle x^{s_2} \rangle_{Id}$ being equivalent.

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